

Revenue-optimal task scheduling and resource management for IoT batch jobs in mobile edge computing

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/Published online: 6March 2020 Received: 31 August 2019 / Accepted: 16 January 2020 •Springer Science+Business Media, LLC, part of Springer Nature 2020

Abstract

(Keywords

1 Introduction 29. 10). \bowtie 34 . A) 39, 102249, , B 180 (B) 2025, 70% , B 102249, , B , B 100876, , B 100101,



16. 3 24 . 16, (). 26 $\lambda -$.B -**5**, **6**. A 16 $\lambda -$ 26 . В 17 . B 2, 3, **5**,

. 6,					22
			,		
2 Related work			32		-
2.1 Peer-to-peer network			,	. 38	
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2.2 Edge computing

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3 System model and problem formulation

3.1 Fundamental notations

System overview

, N M

au , t .

τ.

.

 $\Gamma.$

. $\Gamma \hspace{1cm} \tau. \hspace{1cm} A$ $\Gamma,$

- ,

t

,
$$R_{j}(t)$$
 $R_{j}(t) \leq A_{j}$.

Table 1

t , τ	,
Γ	
i, N, N	, ,
j, M, \mathcal{M}	, , /
B_i	i
L_j	/ j
\mathcal{M}_i	I
N_j	$^{\prime}$
μ_i^{lc}, μ_i^{rc}	\blacksquare ,
$f_i, f_{\mathcal{M}}$	\blacksquare i ,
$ au_i^{lc}, au_i^{rc}$	$\overline{\eta}$, i
δ_i^{rc}	i
w_i, h_i, p_i	, , i
σ	
$d_{i,j}^{tr}, au_{i,j}^{tr}$, i j
τ_i	i
$T_{i,j}$	j i
$A_j, R_j(t)$	j ,
$k, R_j(t)$, j t
$\alpha_{i,j}(t)$	j i t
$\lambda_{i,j}^k(t)$	λ- 26

t. , $j \in \mathcal{M}$. B_i N . A 3.1,

- ,

. 4. $\tau_i \quad ,$

$$i \in \mathcal{N} j \in \mathcal{M}_{i} \quad t = \tau_{i}^{lc} + \tau_{i,j}^{tr} \quad \frac{B_{i}}{\tau_{i}^{rc}} \cdot \alpha_{i,j}(t)$$

$$(4)$$

$$t$$
, $j \in \mathcal{M}$ (5)

$$\alpha_{i,j}(t) \le R_j(t) \qquad \forall j \in \mathcal{M}, \quad \forall t \in \mathcal{T}_{i,j}$$
 (5)

B ,
$$\tau_i$$

$$\alpha_{i,j}(t) = \sum_{i \in \mathcal{N}} \int_{j \in \mathcal{M}_i}^{\tau_i} t = \tau_i^{lc} + \tau_{i,j}^{tr} \frac{B_i}{\tau_i^{rc}} \cdot \alpha_{i,j}(t)$$
(7)

$$\alpha_{i,j}(t) \le R_j(t) \qquad \forall j \in \mathcal{M}, \quad \forall t \in \mathcal{T}_{i,j}$$
 (8)

$$\tau_{i} \qquad \alpha_{i,j}(t) \leq \tau_{i}^{rc} \qquad \forall i \in \mathcal{N} \qquad (9)$$

$$j \in \mathcal{M}_{i} \ t = \tau_{i}^{lc} + \tau_{i,j}^{tr}$$

$$\alpha_{i,j}(t) \in \mathcal{R}_{j}(t) \quad \forall i \in \mathcal{N}, \quad \forall j \in \mathcal{M}_{i}, \quad \forall t \in \mathcal{T}_{i,j} \quad (10)$$

$$\mathcal{R}_{j}(t) \quad \{0, 1, \dots, R_{j}(t)\}, \quad \mathcal{T}_{i,j} \quad \{\tau_{i}^{lc} + \tau_{i,j}^{tr}, \tau_{i}^{lc} + \tau_{i,j}^{tr} + 1, \dots, \tau_{i} \}.$$

(6)

16 ,

, 4.

4 Algorithm design

4.1 Task scheduling framework overview

A 1 - . ,

Γ. Γ. Γ,

 B_i, au_i , etc.) . A , , , , , ,

 α .
(7)
4.2 4.3.

 B_i au_i .

$$\Gamma = \Gamma + 1.$$

4.2 Integral optimum guarantee

A

($\sqrt{}$) 16.

(8)

(9), $C_{r \times c}$ $C_{r \times c}$

$\mathbf{C}_{r \times c}$ 4.3 Equivalent LP transformation (8) (9), . 11 12. **(7)** () r_1 (8) (9). **(7)** $r = r_1 + r_2$ $\alpha_{i,j}(t) \in \mathcal{R}_j(t) =$ $\{0, 1, ..., R_j(t)\},\$ $\tau_{i} - \tau_{i}^{lc} - \tau_{i,j}^{tr} + 1 + N$ $j \in \mathcal{M} i \in \mathcal{N}_{j}$ (11) $\tau_i - \tau_i^{lc} - \tau_{i,j}^{tr} + 1$ $i \in \mathcal{N} \ j \in \mathcal{M}_i$ (12) $\mathbf{C}_{r \times c}$ $\alpha_{i,j}(t)$ 1 $\{0, 1\},\$ (7) -1, 0, 1. $k \in \mathcal{R}_j(t),$ 2, $\lambda_{i,j}^k(t) \in \mathbb{R}^+,$ $\{1, 2, ..., r_1\}$ R_1 . A $R_{j}(t)$ $\alpha_{i,j}(t) = k \cdot \lambda_{i,j}^k(t), \qquad \lambda_{i,j}^k(t) = 1$ $R_1 \quad R_2 = \emptyset.$ y $\lambda_{i,j}^k(t) \in \mathbb{R}^+, \ \forall k \in \mathcal{R}_j(t)$ (8), R_1 1 × c1, . 13. (10), $1 \times c$ R_1 (14). $\alpha_{i,j}(t), \lambda_{i,j}^k(t)$ $i \in \mathcal{N}$ $j \in \mathcal{M}_i$ $t = \tau_i^{lc} + \tau_{i,j}^{tr}$ k = 0 p_i (13)(14) . 15. $a_{x,y} = 1, \ \forall y \in \{1, 2, ..., c\}$ (13) $R_i(t)$ $x \in R_1$ $k \cdot \lambda_{i,j}^{k}(t) \leq R_{j}(t) \quad \forall j \in \mathcal{M}, \quad \forall t \in \mathcal{T}_{i,j}$ (21) $a_{x,y} = 1, \ \forall y \in \{1, 2, ..., c\}$ (14) $x \in R_2$ $\begin{aligned} \tau_i & R_j(t) \\ k \cdot \lambda_{i,j}^k(t) & \leq \tau_i^{rc} & \forall i \in \mathcal{N} \\ j \in \mathcal{M}_i & t = \tau_i^{lc} + \tau_{i,j}^{tr} & k = 0 \end{aligned}$ $a_{x,y} - a_{x,y} = 0 \le 1, \ \forall y \in \{1, 2, ..., c\}$ (15) $x \in R_1$ $x \in R_2$ $\mathbf{C}_{r \times c}$ $\lambda_{i,j}^{k}(t) = 1 \quad \forall i \in \mathcal{N}, \quad \forall j \in \mathcal{M}_{i}, \quad \forall t \in \mathcal{T}_{i,j} \quad (23)$. 8 $\alpha_{i,j}(t), \ \lambda_{i,j}^k(t) \in \mathbb{R}^+ \qquad \forall i \in \mathcal{N}, \qquad \ \forall j \in \mathcal{M}_i,$ $\forall k \in \mathcal{R}_i(t), \quad \forall t \in \mathcal{T}_{i.i}$ 1 (10).**(7)** (20), . 16. 7 ((,(,)- 7 7 (,)

 $\alpha_{i,j}(t) \in 0, R_j(t) \quad \forall i \in \mathcal{N}, \forall j \in \mathcal{M}_i, \forall t \in \mathcal{T}_{i,j}$ (16)

26 .

(17)

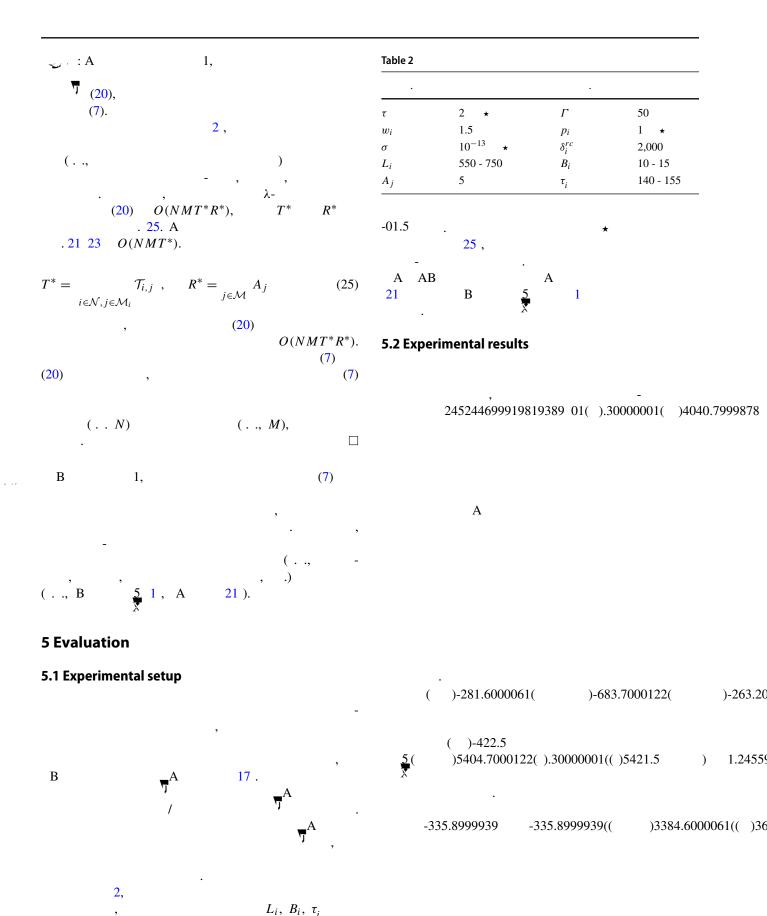
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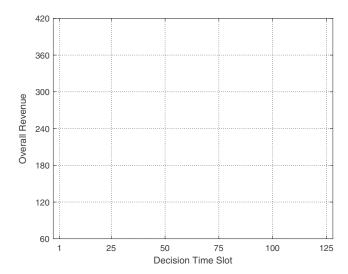
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6 Conclusion and future work

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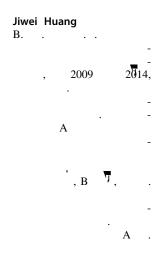
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24.	61 68 (2019)	, , , A 2019 A
25.	270 288	(), B (2016)
26.	34(12):3590 3605 (1977) A	A . A &
27.	15(6):935 7'	946 , A, , <u>5</u> (2019) 12
28.	A, 2	(2010) (2014):466 479 (2010) (
29.	A	, , , , B (2017) - A
30.	, 185. , A,	2016. , , (2019) A
31.	94:351 367	, , , (2019) -
32.	, <u>5</u> ,	57(5):64 69
33.	807 5 , , ,	A 11(4):793 5 (2016) 24(5):2795 2808
34.	$\frac{5}{X}$, , (2)	019) 5 . 3(2):483 493
35.	, ,	5 5 (2010)
36.	56 63	, , , (2019)
37.	, (2019)	6(1):545 556 , A , , ,
38.	:// . /10.1109/	.2019.2901474 , , , (2019) A
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Publisher's note



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